

The construction method of $3X+1$ iterative orbit control and arbitrary step size of orbit

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Abstract: By setting the even - exponent orbit of the $3x + 1$ iteration to construct and solve the initial number, and let the odd - orbit of this initial number iterate under the control of the even - exponent orbit. Through this construction method, it is proved that the $3x + 1$ iteration step length of a specific type of number can be arbitrary or even infinite.

1.Introduction

I . Symbolic Meanings and basic definitions

N^+ is the set of positive integers , R^+ is the set of positive real numbers, $(2, x_0) = 1$ is the co prime of 2, x_0 , Set $x_0, t, n \in N^+, x_0 > 1, (2, x_t) = 1$, and x_0 is odd, We have

$$(1.1) \quad 2^{n_k} x_k = 3x_{k-1} + 1, \quad n_0 = 0$$

It is called the $3x + 1$ conjecture recursion-formula; x_0 is called the initial number (unless otherwise specified, it is an odd number in the following); the step - by - step substitution operation starting from x_0 is called the $3x + 1$ iterative operation (abbreviated as iteration in the following). In the iterative process, the iteration from an odd number x_{k-1} to an even number $3x_{k-1} + 1$ is called an odd - number iteration,

and the iteration from an even number $3x_{k-1}+1$ to an odd number $\frac{3x_{k-1}+1}{2^{n_k}}$ is called an even - number iteration. There are various equation expressions for the $3x + 1$ iterative operation. However, in many cases, it is very convenient to use (1.1) for iterative operations. This article conducts iterative operations based on (1.1), which is helpful for avoiding using heuristic or probabilistic methods to solve problems^[1]. For the convenience of subsequent proofs, we Definition:

Definition 1: denote $T_t(x_0) = \{x_0, x_1, x_2, \dots, x_t\}$, It is called the odd orbit iterated from the attain number x_0 to the cutoff number x_t , Call t the iteration step size

Definition 2: denote $z_t(x_0) = \{n_1, n_2, \dots, n_t\}$, It is called the even exponential orbit with even number iteration from the initial number x_0 to the cutoff number x_t , (Hereinafter referred to as the exponential orbit), if x_k is taken as the initial number, the exponential orbit is denoted as $z_t(x_k) = \{n_{k+1}, n_{k+2}, \dots, n_t\}$.

Definition 3: denote $E_t(x_0) = n_1 + n_2 + \dots + n_t$, $O_t(x_0) = t$, which are respectively called the total number of iterations over even numbers and the total number of iterations over odd numbers from the initial number x_0 up to the number x_t ; denote $D_t(x_0) = E_t(x_0) + O_t(x_0)$, which is called the total number of iterations. If x_k is taken as the initial number, then $E_t(x_k) = n_{k+1} + n_{k+2} + \dots + n_{k+t}$, $O_t(x_k) = t$ etc.

Because it is often used in subsequent studies, we list here a few basic equations. Substituting (1.1) item by item, we have

$$x_k = \frac{1}{2^{n_k}}(3x_0 + 1)$$

$$= \frac{1}{2^{n_k}} \left(\frac{3}{2^{n_{k-1}}} \left(\frac{3}{2^{n_{k-2}}} \left(\dots \left(\frac{3}{2^{n_1}} (3x_0 + 1) + 1 \right) \dots + 1 \right) + 1 \right) + 1 \right)$$

To simplify and combine definitions 3, we have

$$(1.2) \quad 2^{E_k(x_0)} x_k = 3^k x_0 + \sum_{n=0}^{k-1} 3^{k-n-1} \times 2^{E_n(x_0)}, \quad E_0(x_0) = 0, \quad n_0 = 0$$

If x_k is taken as the initial number and x_{k+a} is taken as the terminal number), the above formula can be rewritten as

$$(1.3) \quad 2^{E_a(x_k)} x_{k+a} = 3^a x_k + \sum_{n=0}^{k+a-1} 3^{t+a-n-1} \times 2^{E_n(x_k)}, \quad E_0(x_k) = 0$$

Among them, $E_0(x_k) = 0$ is mainly for the convenience of writing the equation and does not represent the actual value.

II. The $3x + 1$ Conjecture (Collatz Conjecture) and problem-solving idea

1. Conjecture ($3x + 1$ or Collatz Conjecture) : Set $x_0 > 1 \in \mathbb{N}^+$, then x_0 has the total energy $x_k = 1$ through finite step iteration

2. There is an important conjecture in the $3x + 1$ conjecture: the iterative step size of any number is finite .

3. Problem - solving Idea: The difficulty in proving the $3x + 1$ conjecture lies in that the iterative orbit from the initial number x_0 to the terminal number x_t cannot be accurately predicted except through specific

calculations. The problem - solving idea of this paper is that for a given arbitrary exponential orbit $z_t(x) = \{n_1, n_2, \dots, n_t\}$, solve for the initial number x and observe the iterative law of x . This is a brand - new problem - solving tool and provides a new idea for solving the $3x + 1$ conjecture.

III. This paper will prove that:

Constructive theorem: set $h, k \in N^+$, $Z_k(x) = \{1, 1, 1, \dots, 1\}$, where k is any positive integer, then it can always be constructed

$$(1.4) \quad x = (2^{k+1} - 1) + 2^{k+1}h$$

Let the orbit $T_k(x)$ iterate under the control of the exponential orbit $Z_k(x)$ and have

$$(1.5) \quad \lim_{k \rightarrow \infty} x_k = \infty$$

2. Lemma

Lemma 1: Set $a, b, k, m \in N^+$, $2^m a + b | b < 2^m$, $(ab, 2) = 1$, $k < m$, When performing k odd-numbered iterative operations on $2^m a + b$, b respectively, there is

$$(2.1) \quad E_k(b) = E_k(a \cdot 2^m + b) | E_k(b) < m$$

Prove that by mathematical induction, when $k = 1$, b , $a \cdot 2^m + b$ is substituted into (1.1) and combined with the definition 3, we have

$$(2.2) \quad 2^{E_1(b)} x_1 = 3b + 1$$

$$(2.3) \quad 2^{E_1(2^m a + b)} x_1' = 3(2^m a + b) + 1$$

By sorting out (2.3) , we have

$$(2.4) \quad 2^{E_1(a.2^m+b)}(x'_1 - 3.a.2^{m-E_1(a.2^m+b)}) = 3b+1$$

Compared with (2.2)(2.4) , it is obvious that $E_1(b) = E_1(a.2^m+b) | E_1(b) < m$, then the lemma is true.

When $k=t$, in terms of (1.2) associative definition 3, we have

$$(2.5) \quad 2^{E_t(b)} x_t = 3^t b + \sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(b)}$$

$$(2.6) \quad 2^{E_t(2^m a+b)}(x'_t - 3^t.2^{m-E_t(2^m a+b)}.a) = 3^t b + \sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(2^m a+b)}$$

Assuming that the left sides of equations (2.5) and (2.6) are equal, we have

$$(2.7) \quad 3^t b + \sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(2^m a+b)} - 3^t b + \sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(b)} = 0$$

Compare (2.5)(2.6) to the right, we have

$$(2.8) \quad E_t(b) = E_t(a.2^m+b) | E_t(b) < m$$

For $k=t+1$, According to (1.2) and definition 3 we have

$$(2.9) \quad \begin{aligned} 2^{E_{t+1}(b)} x_{t+1} &= 3^{t+1} b + \sum_{n=0}^t 3^{t-n-1} \times 2^{E_n(b)} \\ &= 3^{t+1} b + 2^{E_t(b)} + 3 \left(\sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(b)} \right) \end{aligned}$$

$$(2.10) \quad \begin{aligned} 2^{E_{t+1}(2^m a+b)}(x'_{t+1} - 3^{t+1}.2^{m-E_{t+1}(2^m a+b)}.a) &= 3^{t+1} b + \sum_{n=0}^t 3^{t-n-1} \times 2^{E_n(2^m a+b)} \\ &= 3^{t+1} b + 2^{E_{t+1}(2^m a+b)} + 3 \sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(2^m a+b)} \end{aligned}$$

(2.10)(2.9) minus the left hand side of the equation, combined with

(2.7)(2.8) we have

$$(2.11) \quad \left(2^{E_t(2^m a+b)} - 2^{E_t(b)} \right) + 3 \left(\sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(2^m a+b)} - \sum_{n=0}^{t-1} 3^{t-n-1} \times 2^{E_n(b)} \right) = 0$$

Compare (2.10)(2.9) to the right hand side of the equation

$$(2.12) \quad E_{t+1}(b) = E_{t+1}(a \cdot 2^m + b) \mid E_{t+1}(b) < m$$

Then the lemma is true and the lemma is proved. For other types of proofs of this lemma, see [2].

Lemma 2: Set $x, h \in N^+$, $z_k(x) = \{n_1, n_2, \dots, n_k\}$ be the exponent orbit of the iteration of any positive odd integer x . By choosing an appropriate h , we can always construct

$$(2.13) \quad x = r + 2^{E_{k+1}(x_0)} h \quad \mid 0 < r < 2^{E_{k+1}(x_0)}$$

such that $T_{k+1}(x)$ iterates under the control of $z_k(x)$.

Proof: Set $z_k(x) = \{n_1, n_2, \dots, n_k\}$, According to Definition 3 and (1.2), construct the following equation:

$$(2.14) \quad -r_k = \sum_{n=0}^k 3^{k-n} 2^{E_k(x)}, \quad E_0(x) = 0$$

Do the following for (2.14)

$$(2.15) \quad R \equiv 3^{\varphi(2^{E_{k+1}(x)})-k-1} r_k \equiv 3^{2^{E_{k+1}(x)-1}-k-1} r_k \pmod{2^{E_{k+1}(x)}}$$

Where $\varphi(2^{E_{k+1}(x)}) = 2^{E_{k+1}(x)-1}$ is an Euler function, $0 < R < 2^{E_{k+1}(x)}$.

Set

$$(2.16) \quad r \equiv 2^{E_{k+1}(x)} + R \pmod{2^{E_{k+1}(x)}}$$

Where $0 < r < 2^{E_{k+1}(x)}$, there is always an integer h , such that

$$(2.17) \quad x = r + 2^{E_{k+1}(x)} h$$

Obviously, as long as the appropriate value of h is chosen, the initial number x can be determined, such that $T_k(x) = \{x, x_1, \dots, x_k\}$ iterates under the control of $z_k(x)$, according to the definition 3, $E_{k+1}(x) = E_k(x) + n_{k+1}$, where $n_{k+1} \geq 1$, According to Lemma 3, $E_k(x) + m$, for $m = 1, 2, 3, \dots$ and this can only be determined by substituting term by term and performing trial calculations. Let's illustrate with an example.

Let an arbitrary exponential orbit be $Z_6(x) = \{2, 1, 3, 2, 1, 3\}$. Obviously, $E_1(x) = 2, E_2(x) = 2 + 1 = 3, E_3(x) = 2 + 1 + 3 = 6, \dots, E_6(x) = 12$.

Substituting into (2.14), we have:

$$\begin{aligned} -r_6 &= \sum_{n=0}^{7-1} 3^{k-n-1} 2^{E_k(x)} \\ &= (3^6 + 3^5 \cdot 2^2 + 3^4 \cdot 2^3 + 3^3 \cdot 2^6 + 3^2 \cdot 2^8 + 3 \cdot 2^9 + 2^{12}) \\ &= 12013 \end{aligned}$$

Substituting $E_7(x) = E_6(x) + 1 = 13$ into (2.15), we have

$$R \equiv 3^{2^{E_6(x)} - 7} r_6 \equiv 3^{4089} \times (-12013) \equiv -7527 \pmod{2^{13}}$$

Substituting (2.16) we have

$$r \equiv 2^{E_6(x)+1} + R = 8192 - 7527 \equiv 665 \pmod{2^{13}}$$

The x solution is:

$$(2.18) \quad x_{01} = 665 + 2^{13} h_1$$

By the same token, substituting (2.15) for $E_7 = E_6(x) + 2 = 14$, we have

$$R \equiv 3^{2^{E_6(x)+2} - 7} r_6 \equiv -7527 \pmod{2^{14}}$$

The x solution is:

$$(2.19) \quad x_{02} = 8857 + 2^{14} h_2$$

Take $h_1 = 0,1, h_2 = 0,1$ to verify the following:

Table of Iteration Orbits of Odd Numbers Under the Control of
Exponential Orbits

Tabel1:

		x	x_1	x_2	x_3	x_4	x_5	x_6
$T_6(x_{01})$	$h_1 = 0$	665	499	749	281	211	317	119
	$h_1 = 1$	8857	6643	9965	3737	2805	4205	1577
$Z_6(x_{01})$		0	2	1	3	2	1	3
$T_6(x_{02})$	$h_2 = 0$	8857	6643	9965	3737	2805	4205	1577
	$h_2 = 1$	41625	31219	46829	17561	13171	19757	7409
$Z_6(x_{02})$		0	2	1	3	2	1	3

In this example, when $E_7(x) = E_6(x) + 1 = 13$ satisfies the iteration under the control of $z_6(x)$ orbit, so (2.18) is the minimum solution, and other solutions can be deduced from this minimum solution.

Lemma 4: Set $z_k(x) = \{n_1, n_2, \dots, n_k\}$ be the exponential orbit of the iteration of any positive odd number x , and $T_k(x) = \{x, x_1, x_2, \dots, x_k\} | x_i \neq x_j, 1 \leq i, j \leq k$ be the odd - number orbit generated under the control of $z_k(x)$, If

$$(2.20) \quad \frac{E_m(x)}{m} \leq \frac{E_{m-1}(x)}{m-1} < \frac{\log 3}{\log 2} | 2 \leq m \leq k$$

Then there is always

$$(2.21) \quad x < x_1 < x_2, \dots < x_k$$

Prove by mathematical induction: When $k = 1$, $n_1 = \frac{E_1(x)}{1} < \frac{\log 3}{\log 2}$, x

is the initial number, Substitute into (1.1), we have

$$(2.22) \quad \frac{x_1}{x} = \frac{3}{2^{n_1}} \left(1 + \frac{1}{3x} \right) > 1$$

The lemma is true.

Suppose that the lemma is true when $k = m - 1$, that is,

$$\frac{E_{m-1}(x)}{m-1} < \frac{\log 3}{\log 2}, \quad n_{m-1} < \frac{\log 3}{\log 2} \text{ associated with defining } 3, \text{ we have}$$

$$(2.23) \quad \begin{aligned} n_m = E_m(x) - E_{m-1}(x) &= \frac{E_m(x)}{m} m - \frac{E_{m-1}(x)}{m-1} (m-1) \\ &= \left[\frac{E_m(x)}{m} - \frac{E_{m-1}(x)}{m-1} \right] + \frac{E_{m-1}(x)}{m-1} \end{aligned}$$

substitutes $\frac{E_m(x)}{m} - \frac{E_{m-1}(x)}{m-1} \leq 0$ into (2.23), there is always

$$(2.24) \quad n_m < \frac{\log 3}{\log 2}$$

So, when $k = m$, the lemma is also true. According to (1.1), we have

$$(2.25) \quad \frac{x_m}{x_{m-1}} = \frac{3}{2^{n_m}} \left(1 + \frac{1}{3x_{m-1}} \right)$$

Obviously, when (2.24) is substituted into (2.25), there is always $3/2^{n_m} > 1$, we have

$$(2.26) \quad x_{m-1} < x_m$$

Substituting (1.1) One by one, we have

$$(2.27) \quad x < x_1 < x_2, \dots < x_k$$

End of proof of lemma.

Corollary: Under the control of the exponential orbit $z_k(x) = \{1, 1, \dots, 1\}$ for the odd - orbit $T_k(x) = \{x, x_1, x_2, \dots, x_k\} | x_i \neq x_j, 1 \leq i, j \leq k$, there is

always

$$(2.28) \quad x < x_1 < x_2, \dots < x_k$$

This corollary can be directly deduced from the lemma. Readers can prove it by themselves, and no further elaboration will be given here.

Due to the importance of this corollary in the subsequent proof, we will conduct a calculation verification here. We take $z_9(x) = \{1,1,1,1,1,1,1,1,1\}$ and substitute it into (2.14). Then we have:

$$\begin{aligned} -r_9 &= \left[\begin{array}{l} 3^9 + 3^8 \cdot 2 + 3^7 \cdot 2^2 + 3^6 \cdot 2^3 + 3^5 \cdot 2^4 + 3^4 \cdot 2^5 \\ + 3^3 \cdot 2^6 + 3^2 \cdot 2^7 + 3 \cdot 2^8 + 2^9 \end{array} \right] \\ &= 58025 \end{aligned}$$

Substitute into (2.15):

$$R \equiv 3^{\phi(E_9(x)+1)-k} r_9 \equiv -58025 \times 3^{502} \equiv -1 \pmod{2^{10}}$$

Substitute into (2.16):

$$r \equiv 2^{10} - 1 \equiv 1023 \pmod{2^{10}}$$

Substitute into (2.17):

$$(2.29) \quad x = 1023 + 2^{10} h$$

Taking $h = 0,1,2$ and substituting in (2.29), we have $x = 1023, 2047, 3071$

and substituting in (1.1), the calculations are validated as follows

Table of Iteration Orbits of Odd Numbers Under the Control of
Exponential Orbits

Tabel2:

	x	x_1	x_2	x_3	x_4	x_5	x_6	x_7	x_8	x_9
$T_9(x)$	1023	1353	2303	3455	5183	7775	11663	17495	26243	39365
	2047	3071	4607	6911	10367	15551	23327	34991	52487	78731
	3071	4067	6911	1032	15551	23327	34991	52487	78731	118097
$Z_9(x)$	0	1	1	1	1	1	1	1	1	1

It is proved that the exponential orbits of $T_9(x)$ are $z_9(x) = \{1,1,1,1,1,1,1,1,1\}$ and satisfy $x < x_1 < x_2, \dots < x_9$.

3. Proof of the theorem

Construction theorem: Set $h, k \in N^+$, where k is any positive integer, $Z_k(x) = \{1,1,1, \dots, 1\}$, can always be constructed

$$(3.1) \quad x = (2^{k+1} - 1) + 2^{k+1}h$$

Let the orbit $T_k(x)$ iterate under the control of the exponential orbit $Z_k(x)$, and there is

$$(3.2) \quad \lim_{k \rightarrow \infty} x_k = \infty$$

Proof: in terms of $Z_k(x) = \{1,1,1, \dots, 1\}$, in conjunction with Lemma2 (2.14), always construct

$$(3.3) \quad -r_k = \sum_{n=0}^k 3^{k-n} 2^{E_k(x)} = \frac{3^{k+1} - 2^{k+1}}{3 - 2}$$

Constructed from (2.15) total energy

$$(3.4) \quad R \equiv 3^{\varphi(2^{k+1}) - k - 1} r_k \equiv 3^{2^k - k - 1} r_k \equiv -1 \pmod{2^{k+1}}$$

There $\varphi(2^{k+1}) = 2^k$ is an Euler function, $0 < R < 2^{k+1}$, constructed from (2.16) total energy

$$(3.5) \quad r \equiv 2^{k+1} - 1 \pmod{2^{k+1}}$$

constructed from (2.17) total energy

$$(3.6) \quad x = (2^{k+1} - 1) + 2^{k+1} h$$

Obviously, as long as an appropriate value of h is chosen, the initial number x can be determined, so that $T_k(x) = \{x, x_1, \dots, x_k\}$ iterates under the control of $z_k(x)$. According to Definition 3 $E_{k+1}(x) = E_k(x) + 1$, $n_{k+1} = 1$, and there is always

$$(3.7) \quad \frac{E_m(x)}{m} = \frac{E_{m-1}(x)}{m-1} = 1 \quad | 2 \leq m \leq k$$

According to Lemma 4 and the corollary, at this time, $x < x_1 < x_2 < \dots < x_k$, and obviously there is

$$(3.8) \quad \lim_{k \rightarrow \infty} x_k = \infty$$

Theorem proof, the theorem readers can combine (3.1), lemma 4 inference selfcalculation verification.

4. Several Explanations

(1) The exponential orbit $Z_k(x) = \{1, 1, 1, \dots, 1\}$ in the construction theorem is a form of digital characteristics. There are many such orbits. As long as they meet the conditions of Lemma 4, countless odd - number orbits $T_k(x)$ with arbitrary step lengths can be constructed under the control of the exponential $Z_k(x)$ orbit.

(2) According to (3.5), the odd iterative step size of an odd number $2^{k+1} - 1$ is greater than k . Therefore, for a specific positive integer, regardless of the total iterative step size, the user-defined

iterative step size k can be arbitrarily long or even infinitely large.

(3) Even if we can construct an initial number x with a very large step size, it is still unknown whether $x_k = 1$ can necessarily be achieved through iterations with a finite step size. However, through the construction theorem, we have proved a conjecture in the $3x+1$ conjecture, that is, the problem of "the iterative step size of any positive integer under $3x+1$ iteration is always finite." Obviously, the answer is negative. For specific numbers, there exist many initial numbers with arbitrary step sizes or even infinitely large step sizes.

References

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