

Abstract simulation, basic logic of synchronization algorithm,  
reasoning judgment and hypothesis contradiction

Any even number greater than 2 must be the sum of two  
prime

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Goldbach conjecture

[integer theory]  $2N = P_1 + P_2$

Arithmetic logic judgment theory

Prime number, also known as prime number, has infinite numbers. A natural number greater than 1 can not be integratedly divided by other natural numbers except 1 and itself. In other words, the number has no other factors except 1 and itself; otherwise, it is called composite.

According to the basic theorem of arithmetic, every integer larger than 1 can be written either as a prime number itself or as a product of a series of prime numbers;

furthermore, if the order of these prime numbers in the product is not taken into account, the written form is unique. The smallest prime number is 2.

set up : (natural number , N)  $N > 1$

even number =  $2N$

odd number =  $2N - 1$

$\therefore 2N \div 2 = N$  (Satisfying integer solution)

$\therefore$  natural number  $N > 1$  (There's no prime in an even number.)

refer to: Reference: Euclid's theory of infinity of prime numbers

On the contrary, if the natural number  $2N$  is not equal to the prime number,

On the contrary, prime numbers can only appear in odd numbers.

Or Goldbach's guess is right.  $2N = P_1 + P_2$

Or Goldbach's guess doesn't hold.  $2N \neq P_1 + P_2$

What are the conditions for establishing and what are the conditions for not holding

suppose:  $2N \neq P_1 + P_2$

Equal:  $N \neq P$  ( $P =$  Prime number)

$N =$  even number

N=odd number

$$(\text{even number: } N) - 1 = S_1$$

$$(\text{odd number: } N) - 2 = S_2$$

$$2N - S_1 (\text{N=even number}) = L_1$$

$$2N - S_2 (\text{N=odd number}) = L_1$$

$L_1$  (Either prime or compound.)

Hypothesis:  $L_1$  prime number

$$2N - L_1 = (\text{Either prime or compound.})$$

$$2N - L_1 = (\text{Assuming that it is equal to; } A_1 \times B_1 \times C_1)$$

$$2N - S_1 = L_1 (\text{suppose: } L_1 = \text{composite}$$

number) =  $A_1 \times B_1 \times C_1$  (Simulation. The basic logic of arithmetic.)

Hypothesis :  $L_1 =$  prime number

$$\text{like that: } 2N - L_1 = S_1 (\text{suppose: } S_1 = \text{composite number}) = A_1 \times B_1 \times C_1$$

Extraction prime number、  $A_1 \cdot B_1 \cdot C_1$

$$\left\{ \begin{array}{l} 2N - A_1 = L_2 (\text{Integer decomposition}) = A_2^n \\ 2N - B_1 = H_2 (\text{Integer decomposition}) = B_2^n \\ 2N - C_1 = M_2 (\text{Integer decomposition}) = C_2 \end{array} \right.$$

$L_2, H_2, M_2,$  (Either prime or compound.)

Hypothesis;  $L_2, H_2, M_2,$  It's a compound number、

Extraction prime number:  $A_2, B_2, C_2,$

$$\left\{ \begin{array}{l} 2N - A_2 = L_3 (\text{Integer decomposition}) = A_3^n \end{array} \right.$$

$$2N - B_2 = H_3(\text{Integer decomposition}) = B_3^n$$

$$2N - C_2 = M_3(\text{Integer decomposition}) = C_3^n$$

$L_3$ 、 $H_3$ 、 $M_3$ 、 (Either prime or compound.)

Hypothesis;  $L_3$ 、 $H_3$ 、 $M_3$ 、 It's a compound number.

Extraction prime number:  $A_3$ 、 $B_3$ 、 $C_3$ 、

$$\left\{ \begin{array}{l} 2N - A_3 = L_4(\text{Integer decomposition}) = A_4^n \\ 2N - B_3 = H_4(\text{Integer decomposition}) = B_4^n \\ 2N - C_3 = M_4(\text{Integer decomposition}) = C_4^n \end{array} \right.$$

$L_4$ 、 $H_4$ 、 $M_4$ 、 (Either prime or compound.)

Hypothesis;  $L_4$ 、 $H_4$ 、 $M_4$ 、 It's a compound number.

Extraction prime number:  $A_4$ 、 $B_4$ 、 $C_4$ 、

$$\left\{ \begin{array}{l} 2N - A_4 = L_5(\text{Integer decomposition}) = A_5^n \\ 2N - B_4 = H_5(\text{Integer decomposition}) = B_5^n \\ 2N - C_4 = M_5(\text{Integer decomposition}) = C_5^n \end{array} \right.$$

$L_5$ 、 $H_5$ 、 $M_5$  (Either prime or compound.)

Hypothesis:  $L_5$ 、 $H_5$ 、 $M_5$ 、 It's a compound number.

Extraction prime number:  $A_5$ 、 $B_5$ 、 $C_5$ 、

.....Analog arithmetic logic ( WY1 式 )

(一) 、 Either arithmetic logic, finite loop (prime cycle)

The prime cycle stands for  $2N$ . Using this arithmetic logic, we judge

that we are all compound numbers. ( $2N \neq P_1 + P_2$ )

Simulation: arithmetic, finite cycle hypothesis for all complex numbers

$$2N - A = B^b \text{ (Extraction prime number: } B \text{)}$$

$$2N - B = C^c \text{ (Extraction prime number: } C \text{)}$$

$$2N - C = A^a \text{ (Extraction prime number: } A \text{)}$$

(二)、 Either the arithmetic logic is infinitely noncyclic (arithmetic prime is not cyclic) (infinite increment of different prime numbers)

The prime number is not cyclic, which means  $2N$ . Using this arithmetic logic, it's all a compound number. ( $2N \neq P1 + P2$ )

The representation of an infinitely noncyclic prime number (infinitely increasing)

Any  $2N$  condition is finite

$$2N - P \text{ (Infinite number of different primes: } 2 < P < 2N \text{)}$$

Finite and infinite contradictions、 Vice versa by the above logic ( $2N = P1 + P2$ )

Hypothesis: ( $2N \neq P1 + P2$ )

extract (一) Finite arithmetic cycle、 ( $2N \neq P1 + P2$ )

And then do the analog arithmetic judgment.

set up: ( $2N \neq P1 + P2$ )

Here,  $E_2$ 、  $F_2$ 、  $G_2$ 、 It can be prime or compound.

$$\textcircled{2} \left\{ \begin{array}{l} 2N - 2A_2 = 2E_2 \div 2(E_2, \text{ Integer decomposition}) = S_2^n \\ 2N - 2B_2 = 2F_2 \div 2(F_2, \text{ Integer decomposition}) = W_2^n \end{array} \right.$$

$$2N - 2C_2 = 2G_2 \div 2(G_2, \text{ Integer decomposition } ) = R_2^n$$

Extraction prime number:  $S_2, W_2, R_2$

$$\textcircled{1} \left\{ \begin{array}{l} 2N - S_2 = E_3(E_2, \text{ Integer decomposition } ) = S_3^n \\ 2N - W_2 = F_3(F_2, \text{ Integer decomposition } ) = W_3^n \\ 2N - R_2 = G_3(G_2, \text{ Integer decomposition } ) = R_3^n \end{array} \right.$$

$E_3, F_3, G_3$ , ( Either prime or compound. )

Hypothesis ;  $E_3, F_3, G_3$ , It's a compound number.

Extraction prime number:  $S_3, W_3, R_3$

$$\textcircled{2} \left\{ \begin{array}{l} 2N - 2S_3 = 2E_4 \div 2(E_4, \text{ Integer decomposition } ) = S_4^n \\ 2N - 2W_3 = 2F_4 \div 2(F_4, \text{ Integer decomposition } ) = W_4^n \\ 2N - 2R_3 = 2G_4 \div 2(G_4, \text{ Integer decomposition } ) = R_4^n \end{array} \right.$$

Here,  $E_4, F_4, G_4$ , It can be prime or compound.

Extraction prime number:  $S_4, W_4, R_4$ ,

$$\textcircled{1} \left\{ \begin{array}{l} 2N - S_4 = E_5(E_5, \text{ Integer decomposition } ) = S_5^n \\ 2N - W_4 = F_5(F_5, \text{ Integer decomposition } ) = W_5^n \\ 2N - R_4 = G_5(G_4, \text{ Integer decomposition } ) = R_5^n \end{array} \right.$$

$E_5, F_5, G_5$ , ( Either prime or compound. )

Hypothesis ;  $E_5, F_5, G_5$ , It's a compound number.

Extraction prime number:  $S_5, W_5, R_5$ ,

$$\left\{ \begin{array}{l} 2N - 2S_5 = 2E_6 \div 2(E_6, \text{ Integer decomposition } ) = S_6^n \end{array} \right.$$

$$\textcircled{2} \quad 2N - 2W_5 = 2F_6 \div 2(F_6, \text{ Integer decomposition } ) = W_6^n$$

$$2N - 2R_5 = 2G_6 \div 2(G_6, \text{ Integer decomposition } ) = R_6^n$$

Here,  $E_6$ 、  $F_6$ 、  $G_6$ 、 It can be prime or compound. Extraction prime number、  $S_6$ 、  $W_6$ 、  $R_6$ 、 .....

Analog arithmetic basic logic ( WY2 式 )

( 三 ) 、 Or finite arithmetic cycle hypothesis (prime cycle)

Finite arithmetical cyclic logic, representing  $2N$  that uses this arithmetic logic to judge that all are complex numbers ( $2N \neq P_1 + P_2$ )

( 四 ) 、 Arithmetic no cycle (Primes do not circulate) (infinitely increasing)

The prime number is not cyclic, which means  $2N$ . Using this arithmetic logic, it's all a compound number. ( $2N \neq P_1 + P_2$ )

(infinite increment of different prime numbers)

$2N - P$  (Infinite number of different primes:  $P < 2N$ )

The contradiction between finite and infinite; vice versa ( $2N = P_1 + P_2$ )

Hypothesis : ( $2N \neq P_1 + P_2$ ) Only choice. (一) 、 (三) Arithmetic finite cycle (prime cycle)

Two arithmetic circle logic problems are merged into one problem, reasoning and judgment. ( WY2 ) Or you get prime numbers and ( WY1 ) All the same

( WY2 ) Or you get prime numbers and , ( WY1 ) Existence

of Incrementally Different Primes

set up : ( WY2 ) All cyclic prime numbers, Same as ( WY1 ) prime number.

set up : ( WY1 ) Compound arithmetic cyclic logic S item column

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$

abstract : S Arithmetic loop simulation:

$$2N - A = B^b$$

$$2N - B = C^c$$

$$2N - C = D^d$$

$$2N - D = E^e$$

$$2N - E = F^f$$

$$2N - F = G^g$$

$$2N - G = H^h$$

$$2N - H = A^a$$

Then all the solutions above are compound numbers. Solution :  $2N$  is not equal to the sum of two prime numbers

( WY2 ) Compound arithmetic cyclic logic reference , refer to

( WY1 ) prime number

$$(2N - A = B^b) \quad (2N - 2A = 2B^d)$$

$$2N - A = B^b \text{ 和 } 2N - 2A = 2B^d$$

$$A = B^b - 2B^d$$



$$A = B^d(B^{b-d} - 2)$$

$$2N - B^d(B^{b-d} - 2) = B^b$$

$$2N - B^b + 2B^d = B^b$$

$$N = 2B^b - 2B^d$$

B is not a prime factor of 2N

$2N \div B < \text{Unsatisfied integer solution} >$

Hypothesis: Equality holds (fractional solution)

The contradiction between integer theory and hypothesis

② A → Derivation of Prime numbers without B

$$(2N - 2A = 2B^d) \text{ false}$$

Here (WY2) and (WY1) get the same prime number

hypothesis:  $2N - 2A = 2C$  (Extraction prime number : C)

① Arithmetic logic derivation

$$2N - C = D^d \text{ (Extraction prime number : D)}$$

get into :  $2N - 2D =$

If it's all the same, then it must be. ② A → reduction C, D, , , , ,

If: C ① → Derivation D

Hypothesis: C ① → Derivation is not possible. D

like that , ( WY1 式 ) C → reduction D

The same even number minus the same prime number

there are different prime numbers

$$2N - C = D^d$$

$$2N - C = E^E$$

If the hypothesis :  $D^d = E^E$

like that:  $D^d \div E$  ( Unsatisfied integer solution )

Contradiction with propositional integer theory

Conversely, the assumption is not true

Vice versa (WY1) and (WY2) must produce different prime numbers

Arithmetic cycle simulation:  $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$

Abstract hypothesis: (WY2) and (WY1) arithmetic Prime cyclic

Simulation

$$\textcircled{2}B \rightarrow \textcircled{1}E \rightarrow \textcircled{2}F \rightarrow \textcircled{1}A \rightarrow B \textcircled{2}$$

Simulated judgment (WY1) and (WY2) relation

Simulation (WY1) finite arithmetic cycle hypothesis

$$2N - A = B^b$$

$$2N - B = C^c$$

$$2N - C = D^d$$

$$2N - D = E^e$$

$$2N - E = F^f$$

$$2N - F = G^g$$

$$2N - G = H^h$$

$$2N - H = A^a$$

$$2N - A = B^b$$

substitution

$$2N - E^e = D$$

$$2N - C = (2N - E^e)^d$$

$$C = 2N - (2N - E^e)^d$$

$$2N - B = [2N - (2N - E^e)^d]^c$$

$$B = 2N - [2N - (2N - E^e)^d]^c$$

代入

$$2N - 2B = 2E^M$$

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^M$$

So there's only the relationship between 2N and E here.

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^M$$

Hypothesis : P and K are equal to any integer

$$2N - 2 [2NP - KE^x] = 2E^M$$

$$N - [2NP - KE^x] = E^M$$

$$N - 2NP + KE^x = E^M$$

$$N(1 - 2P) = E^M - KE^x$$

$$N(1 - 2P) = E^x(E^{M-x} - K)$$

And here we have the relationship between the existence of prime factors of N and E

The prime factor of N does not include E

If the equation holds  $P$  and  $K$ , there is a non-integer solution.

It is assumed that there is a contradiction between the theory of fraction and the theory of propositional integers. On the contrary, the assumption  $(2N \neq P_1 + P_2)$  can only choose (one), (three) finite hypothesis (prime cycle) to exist, and the finite hypothesis does not hold.

(一)、(三) There is an infinite arithmetic non-cyclic hypothesis  
(Existence of infinitely increasing different prime numbers)

$2N$  finite value and infinite contradiction

Vice versa by the above arithmetic logic  $(2N \neq P_1 + P_2)$  false

So, by the above arithmetic logic,  $(2N = P_1 + P_2)$

References:

Euclidean prime number infinite

The problem of Cantor's Continuum Cardinal number.

In 1874, Cantor conjectured that there was no other cardinality between the countable set cardinality and the real set cardinality, the famous continuum hypothesis. Godel, an Austrian mathematical logicist living in the United States, proved that there is no contradiction between the continuum

hypothesis and the ZF set theory axiom system. In 1963, P.Choen, an American mathematician, proved that the continuum hypothesis and the ZF axiom were independent of each other. Therefore, the continuum hypothesis cannot be proved by ZF axiom. In this sense, the problem has been solved.

(2) the noncontradiction of arithmetic axiom system.2)

Distinguished Mr. (Madam)

I'm sorry.

My educational level is limited, the paper is the end part, need relevant mathematical theory to revise and write more perfect. This paper is based on the existence of an infinite number of even numbers.

抽象模拟，同步算术基本逻辑，推理判断和假设矛盾

大于 2 任意偶数必定是两个素数之和

哥德巴赫猜想

【整数理论】 $2N=P_1+P_2$

Arithmetic logic judgment theory

哥德巴赫猜想，命题大于 2 的偶数等于两个素数之和。

那么自然数，素数，

命题可以我们就可以判断属于【整数理论】

数(prime number)又称素数,有无限个。一个大于 1 的自然数,除了 1 和它本身外,不能被其他自然数整除,换句话说就是该数除了 1 和它本身以外不再其他的因数;否则称为合数。

根据算术基本定理,每一个比 1 大的整数,要么本身是一个质数,要么可以写成一系列质数的乘积;而且如果不考虑这些质数在乘积中的顺序,那么写出来的形式是唯一的。最小的质数是 2

设: (自然数  $N$ )  $N > 1$

偶数  $= 2N$

奇数  $= 2N - 1$

$\therefore 2N \div 2 = N$  (满足整数解)

$\therefore$  自然数  $N > 1$  (偶数里面没有素数)

参照: 欧几里得素数无穷大理论

反之当自然数  $2N$  不等于素数

反之素数只能出现在奇数里面

要么哥德巴赫猜想正确、 $2N = P_1 + P_2$

要么哥德巴赫猜想不成立  $2N \neq P_1 + P_2$

成立的条件是什么、不成立的条件是什么

假设:  $2N \neq P_1 + P_2$

当  $N \neq P$  ( $P$  为素数)

$N = \text{偶数}$

$N = \text{奇数}$

(偶数:  $N$ )  $- 1 = S_1$

(奇数:  $N$ )  $- 2 = S_2$

$2N - S_1$  ( $N = \text{偶数}$ )  $= L_1$

$2N - S_2$  ( $N = \text{奇数}$ )  $= L_1$

$L_1$  (要么是素数、要么是复合数)

假设:  $L_1$  素数

$2N - L_1 =$  (要么是素数、要么是复合数)

$2N - L_1 =$  (假设等于  $A_1 \times B_1 \times C_1$ )

$2N - S_1 = L_1$  (假设:  $L_1 = \text{复合数}$ )  $= A_1 \times B_1 \times C_1$  (模拟. 算术基本逻辑.)

如果  $L_1 = \text{素数}$ 。

那么  $2N - L_1 = S_1$  (假设:  $S_1 = \text{复合数}$ )  $= A_1 \times B_1 \times C_1$

抽取素数、 $A_1, B_1, C_1$

$$\left\{ \begin{array}{l} 2N - A_1 = L_2 (\text{整数分解}) = A_2^n \\ 2N - B_1 = H_2 (\text{整数分解}) = B_2^n \\ 2N - C_1 = M_2 (\text{整数分解}) = C_2^n \end{array} \right.$$

$L_2, H_2, M_2$ 、(要么是素数、要么是复合数)

假设:  $L_2, H_2, M_2$ 、是复合数、

抽取素数：  $A_2$ 、  $B_2$ 、  $C_2$ 、

$$\left\{ \begin{array}{l} 2N - A_2 = L_3(\text{整数分解}) = A_3^n \\ 2N - B_2 = H_3(\text{整数分解}) = B_3^n \\ 2N - C_2 = M_3(\text{整数分解}) = C_3^n \end{array} \right.$$

$L_3$ 、  $H_3$ 、  $M_3$ 、 （要么是素数、要么是复合数）

假设：  $L_3$ 、  $H_3$ 、  $M_3$ 、 是复合数

抽取素数：  $A_3$ 、  $B_3$ 、  $C_3$ 、

$$\left\{ \begin{array}{l} 2N - A_3 = L_4(\text{整数分解}) = A_4^n \\ 2N - B_3 = H_4(\text{整数分解}) = B_4^n \\ 2N - C_3 = M_4(\text{整数分解}) = C_4^n \end{array} \right.$$

$L_4$ 、  $H_4$ 、  $M_4$ 、 （要么是素数、要么是复合数）

假设：  $L_4$ 、  $H_4$ 、  $M_4$ 、 是复合数

抽取素数：  $A_4$ 、  $B_4$ 、  $C_4$ 、

$$\left\{ \begin{array}{l} 2N - A_4 = L_5(\text{整数分解}) = A_5^n \\ 2N - B_4 = H_5(\text{整数分解}) = B_5^n \\ 2N - C_4 = M_5(\text{整数分解}) = C_5^n \end{array} \right.$$

$L_5$ 、  $H_5$ 、  $M_5$  （要么是素数、要么是复合数）

假设：  $L_5$ 、  $H_5$ 、  $M_5$ 、 是复合数

抽取素数：  $A_5$ 、  $B_5$ 、  $C_5$ 、

.....模拟算术逻辑（WY1 式）

（一）、 要么、 算术逻辑有限循环（素数循环）



素数循环代表  $2N$  用这个算术逻辑判断都是复合数 ( $2N \neq P_1 + P_2$ )

模拟：算术有限循环假设 (全部是复合数)

$$2N - A = B^b \text{ (抽取素数: } B)$$

$$2N - B = C^c \text{ (抽取素数: } C)$$

$$2N - C = A^a \text{ (抽取素数: } A)$$

(二)、要么、算术逻辑无限不循环(算术素数不循环) (无限递增不相同素数)

素数不循环代表  $2N$  用这个算术逻辑判断都是复合数 ( $2N \neq P_1 + P_2$ )

不循环无穷的素数代表 (无限递增)

任意  $2N$  条件是有限减  $1=0$

$$2N - P \text{ (无限不相同素数 } 2 < P < 2N)$$

有限和无限矛盾、反之由上面逻辑 ( $2N = P_1 + P_2$ )

假设 ( $2N \neq P_1 + P_2$ )

抽取 (一) 有限算术循环、 ( $2N \neq P_1 + P_2$ )

再进行模拟算术判断

设 ( $2N \neq P_1 + P_2$ )

这里  $E_2$ 、 $F_2$ 、 $G_2$ 、可以是素数或者是复合数

$$\textcircled{2} \left\{ \begin{array}{l} 2N - 2A_2 = 2E_2 \div 2(E_2 \text{ 因式分解}) = S_2^n \\ 2N - 2B_2 = 2F_2 \div 2(F_2 \text{ 因式分解}) = W_2^n \\ 2N - 2C_2 = 2G_2 \div 2(G_2 \text{ 因式分解}) = R_2^n \end{array} \right.$$

抽取素数:  $S_2$ 、 $W_2$ 、 $R_2$

$$\textcircled{1} \left\{ \begin{array}{l} 2N - S_2 = E_3(E_2 \text{ 整数分解}) = S_3^n \\ 2N - W_2 = F_3(F_2 \text{ 整数分解}) = W_3^n \\ 2N - R_2 = G_3(G_2 \text{ 整数分解}) = R_3^n \end{array} \right.$$

$E_3$ 、 $F_3$ 、 $G_3$ 、（要么是素数、要么是复合数）

假设； $E_3$ 、 $F_3$ 、 $G_3$ 、是复合数

抽取素数： $S_3$ 、 $W_3$ 、 $R_3$

$$\textcircled{2} \left\{ \begin{array}{l} 2N - 2S_3 = 2E_4 \div 2(E_4 \text{ 因式分解}) = S_4^n \\ 2N - 2W_3 = 2F_4 \div 2(F_4 \text{ 因式分解}) = W_4^n \\ 2N - 2R_3 = 2G_4 \div 2(G_4 \text{ 因式分解}) = R_4^n \end{array} \right.$$

这里  $E_4$ 、 $F_4$ 、 $G_4$ 、可以是素数或者复合数

抽取素数： $S_4$ 、 $W_4$ 、 $R_4$ 、

$$\textcircled{1} \left\{ \begin{array}{l} 2N - S_4 = E_5(E_5 \text{ 整数分解}) = S_5^n \\ 2N - W_4 = F_5(F_5 \text{ 整数分解}) = W_5^n \\ 2N - R_4 = G_5(G_2 \text{ 整数分解}) = R_5^n \end{array} \right.$$

$E_5$ 、 $F_5$ 、 $G_5$ 、（要么是素数、要么是复合数）

假设： $E_5$ 、 $F_5$ 、 $G_5$ 、是复合数

抽取素数： $S_5$ 、 $W_5$ 、 $R_5$ 、

$$\textcircled{2} \left\{ \begin{array}{l} 2N - 2S_5 = 2E_6 \div 2(E_6 \text{ 因式分解}) = S_6^n \\ 2N - 2W_5 = 2F_6 \div 2(F_6 \text{ 因式分解}) = W_6^n \\ 2N - 2R_5 = 2G_6 \div 2(G_6 \text{ 因式分解}) = R_6^n \end{array} \right.$$

这里  $E_6$ 、 $F_6$ 、 $G_6$ 、可以是素数或者是复合数

抽取素因数  $S_6$ 、 $W_6$ 、 $R_6$ 、

.....模拟算术基本逻辑 (WY2 式)

(三)、要么,有限算术循环(素数循环)

有限算术循环逻辑,代表  $2N$  用这个算术逻辑判断都是复合数

( $2N \neq P_1 + P_2$ )

(四)、要么无限算术不循环假设(素数不循环)(无限递增)素

数不循环代表  $2N$  用这个算术逻辑判断都是复合数( $2N \neq P_1 + P_2$ )

不循环无穷的素数代表(无限递增)

$2N - P$  (无限不相同素数  $P < 2N$ )

有限与无限矛盾由上面逻辑( $2N \neq P_1 + P_2$ )不成立

反之( $2N = P_1 + P_2$ )

假设( $2N \neq P_1 + P_2$ )只能选择(一)、(三)有限算术循环(素数循环)

将两个算术逻辑假设问题合并成一个问题,推理判断。

要么、(WY2)全部素数与(WY1)全部相同

要么、(WY2)全部素数与(WY1)存在递增不相同素数

设:(WY2式)(素数循环)所有素数、与(WY1式)素数全部相同

设:(WY1)复合数算术循环逻辑S项

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$

抽象:S项算术循环模拟:

$$2N - A = B^b$$

$$2N - B = C^c$$

$$2N - C = D^d$$

$$2N - D = E^e$$

$$2N - E = F^f$$

$$2N - F = G^g$$

$$2N - G = H^h$$

$$2N - H = A^a$$

那么上面解的都是复合数

解得：2N 不等于两个素数之和

( WY2 ) 复合数算术循环逻辑参照 ( WY1 )

$$( 2N - A = B^b ) \quad ( 2N - 2A = 2B^d )$$

$$2N - A = B^b \text{ 和 } 2N - 2A = 2B^d$$

$$A = B^b - 2B^d$$

$$A = B^d(B^{b-d} - 2)$$

$$2N - B^d(B^{b-d} - 2) = B^b$$

$$2N - B^b + 2B^d = B^b$$

$$N = 2B^b - 2B^d$$

B 不是 2N 的素因数

$2N \div B$  <不满足整数解>

假设：等式成立（有分数解）

整数理论和假设矛盾

②A→推导得不到 B 的素数

( $2N - 2A = 2B^d$ ) 不成立

这里 (WY2 式) 和 (WY1 式) 得到的素数相同

假设:  $2N - 2A = 2C$  (抽取素数: C)

进入①算术逻辑推导

$2N - C = D^d$  (抽取素数: D)

进入  $2N - 2D =$

如果全部相同那么必定是②A→推导 C,D, , , , ,

如果 C①→推导到 D

假设 C①→推导不到 D

那么 (WY1 式) C→推导 D

相同的偶数减去相同的素数存在不相同素数

$2N - C = D^d$

$2N - C = E^E$

如果假设成立  $D^d = E^E$

那么  $D^d \div E$  没有整数解与命题整数理论存在矛盾

反之假设不成立

反之 (WY1 式) 与 (WY2) 必定产生不相同素数

循环模拟  $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$

抽象假设: (WY2)和(WY1)素数循环模拟成立

$$\textcircled{2}B \rightarrow \textcircled{1}E \rightarrow \textcircled{2}F \rightarrow \textcircled{1}A \rightarrow B \textcircled{2}$$

模拟判断 ( WY1 式 ) 与 ( WY2 式 ) 关系

模拟 ( WY1 式 ) 有限算术循环假设

$$2N - A = B^b$$

$$2N - B = C^c$$

$$2N - C = D^d$$

$$2N - D = E^e$$

$$2N - E = F^f$$

$$2N - F = G^g$$

$$2N - G = H^h$$

$$2N - H = A^a$$

$$2N - A = B^b$$

代入

$$2N - E^e = D$$

$$2N - C = (2N - E^e)^d$$

$$C = 2N - (2N - E^e)^d$$

$$2N - B = [2N - (2N - E^e)^d]^c$$

$$B = 2N - [2N - (2N - E^e)^d]^c$$

代入

$$2N - 2B = 2E^M$$

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^M$$

那么这里只存在关系式  $2N$  与  $E$  的关系式

$$2N - 2[2N - (2N - E^e)^d]^c = 2E^m$$

假设： $P$  和  $K$  等于任意整数)

$$2N - 2[2NP - KE^x] = 2E^M$$

$$N - [2NP - KE^x] = E^M$$

$$N - 2NP + KE^x = E^M$$

$$N(1 - 2P) = E^M - KE^x$$

$$N(1 - 2P) = E^x(E^{M-x} - K)$$

这里就产生  $N$  与  $E$  存在素因数关系

而  $N$  的素因数不包含  $E$

如果等式成立  $P$  和  $K$  存在非整数解

假设存在分数和命题整数论之间存在矛盾。

反之假设 ( $2N \neq P_1 + P_2$ ) 只能选择 (一)、(三) 有限假设 (素数循环) 存在一式有限假设不成立

那么 (一)、(三) 里面两式存在一式为无穷假设 (存在无限递增不相同素数)

$2N$  是有限值与无限不相同素数存在矛盾

反之由上面算术逻辑 ( $2N \neq P_1 + P_2$ ) 不成立

那么由上面算术逻辑 ( $2N = P_1 + P_2$ )

参考文献:

欧几里得反证法

康托的连续统基数问题。

1874 年，康托猜测在可数集基数和实数集基数之间没有别的基数，即著名的连续统假设。1938 年，侨居美国的奥地利数理逻辑学家哥德尔证明连续统假设与 ZF 集合论公理系统的无矛盾性。1963 年，美国数学家科恩 (P.Cohen) 证明连续统假设与 ZF 公理彼此独立。因而，连续统假设不能用 ZF 公理加以证明。在这个意义下，问题已获解决。

(2) 算术公理系统的无矛盾性。

尊敬的先生 (女士)

本人文化水平有限，论文就是结尾部分，需要相关数学理论着进行修改写得更完善一点。(我可以说但是写的不完整)  
论文是以无限大的偶数是否存在反列为基础