

# NUMERICAL ANALYSIS ON CHEN'S AUTHENTICATION SCHEME FOR MULTI-SERVER ENVIRONMENT

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## ABSTRACT

With the fast advancement in the field of network security everything gets the chance to be possible on web. Remote user authentication is an imperative system in the networks framework to check the exactitude of remote user over the public channel. In this authentication procedure, server checks accreditation of the user that user is authentic and legal one or not. For that Server and user commonly confirm each other and make a same session key for encryption of upcoming conversations. There are two types of authentication: Single server and Multi server. To overcome the drawback of single server authentication (remembering id and pswd for accessing each of the server), the concept of Multi server comes, in which user first register with RC, and whatever servers are registered under RC can be accessed by user by providing single id and pswd for all. Here We review US patent [US 9264425 B1] scheme which is based on Multi server authentication, we provide mathematical analysis of the same with some attacks found on it.

*Index Terms*— Remote User Authentication, Multi Server authentication, Network Security, Smart Card, Numerical analysis

## INTRODUCTION

Advances in network communication improved the quality of online available services. Nowadays people prefer online services as it is easy to access without moving from your place and it provides faster accessibility. So the data transfer from user to server and server to user over the public channel and the possibilities of eavesdropping information is more in public channel compared to secure channel. In public channel there are chances of intercept, modify and delete the messages by an attacker. Authentication plans utilizing smart card have increased huge consideration due their appropriateness and ease of use in multi-server environment. So we can also say, by using smart card along with password is helped us to verify the legitimacy of the user.

Actually, an authentication procedure is for checking the correctness of both the user and the server. That simply means, authentication procedure involves mutual authentication between server and user and session key establishment for upcoming conversations. In Multi server environment there are mainly three entities involved: User, Server and RC (Registration Center). In Multi server authentication, user first have to register him/herself with RC, then after user can access any of the server associated with RC by providing same user id and password. So the benefit of this authentication is user doesn't have to remember all the password for accessing all the different service providing server and it is convenient for user.

There are many schemes available which are using smart card. It can be classified into two categories: Two factor and Three factor authentication. In two factor authentication [1-4], password and

smart card are used and in three factor authentication [5-7], password, smart card and biometrics are used. In three factor authentication Biometric properties (fingerprint, retina, face etc.) are involved. The benefit of using biometric properties are like it can't lost, not shared easily, contains unique feature, difficult to guess.

In 2005, Lee et al [1] proposed improved authentication scheme using smart card, which is of two factor authentication. In 2009, Xu et al [2] proposed a scheme in which they assumed that both password and smart card are not stolen at the same time, so they proposed a scheme in which either password or smart card is unknown to the attacker, so attacker does not impersonate the legal user. In 2014, He et al [5] proposed a scheme based on ECC, his scheme overcomes the drawbacks of Yoon et al [7]'s scheme and claimed that his scheme is actually the first correct three factor scheme. In 2015, Li et al [6] proposed a scheme with fuzzy extractor and make comparative study of Yoon et al [7] and Shen et al's scheme.

## 2. NUMERICAL ANALYSIS OF CHEN'S SCHEME

This Patent [US 9264425 B1] [8] is based on two factor authentication, that means it focuses on password and smart card in Multi Server Environment. This patent explains four various procedures like 1) Registration Procedure, 2) Login Procedure, 3) Verification Procedure and 4) Password Changing Procedure. All the procedures are explained here with numerical calculations. Instead of original hash value, we have assumed hash value for simplicity.

- **Assumed Parameters**

User $iD_i = 1$	Password $PW_i = 2$	Random No $r = 3$
Master Key $x = 4$	Secret no. $y = 5$	User Nonce $N_i = 6$
Remote server Identity $SID_j = 7$	Remote server Nonce $N_j = 8$	New Password $PW_{i\text{new}} = 9$
New Random No $r_{\text{new}} = 10$		

- **Assumed Hash Values**

$h(1) = 21$	$h(21) = 22$	$h(4  5) = 23$	$h(22  23) = 24$
$h(1  4) = 25$	$h(1  21) = 36$	$h(25) = 27$	$h(5) = 8$
$h(28  6  7) = 29$	$h(25  24  6) = 37$	$h(24  6  7) = 38$	$h(13  24  6) = 32$
$h(13  6  24  7) = 33$	$h(3) = 36$	$h(1  36) = 39$	$h(13  8  24  7) = 34$
$h(13  8  6  24  7) = 35$			

### Registration Procedure:

$h(r \text{ XOR } PW_i) = h(3 \text{ XOR } 2) = h(1) = 21$   
 $R_i = h(h(r \text{ XOR } PW_i)) = h(h(1)) = h(21) = 22$   
 $M_i = h(R_i || h(x||y)) = h(22 || h(4||5)) = h(22||23) = 24$   
 $E_i = M_i \text{ XOR } h(r \text{ XOR } PW_i) = 24 \text{ XOR } 21 = 13$   
 $L_i = h(ID_i || x) = h(1||4) = 25$   
 $W_i = L_i \text{ XOR } h(ID_i || h(r \text{ XOR } PW_i)) = 25 \text{ XOR } h(1||21) = 25 \text{ XOR } 36 = 61$   
 $F_i = h(L_i) = h(25) = 27$

### Login Procedure:

$L_i = W_i \text{ XOR } h(ID_i || h(r \text{ XOR } PW_i)) = 61 \text{ XOR } h(1||21) = 61 \text{ XOR } 36 = 25$

$F_i^* = h(L_i) = h(25) = 27$  and  $F_i^* = F_i \rightarrow 27=27 \rightarrow$  Yes User is authenticated

$M_i = E_i \text{ XOR } h(r \text{ XOR } PWi) = 13 \text{ XOR } 21 = 24$

$R_i = h(h(r \text{ XOR } PWi)) = h(21) = 22$

$G_{ij} = R_i \text{ XOR } h(h(y) || Ni || SID_j) = 22 \text{ XOR } h(28 || 6 || 7) = 22 \text{ XOR } 29 = 11$

$CID_i = h(r \text{ XOR } PWi) \text{ XOR } h(L_i || M_i || Ni) = 21 \text{ XOR } h(25 || 24 || 6) = 21 \text{ XOR } 37 = 48$

$H_{ij} = L_i \text{ XOR } h(M_i || Ni || SID_j) = 25 \text{ XOR } h(24 || 6 || 7) = 25 \text{ XOR } 38 = 63$

$Z_i = h(E_i || M_i || Ni) = h(13 || 24 || 6) = 32$

First Parameter Set  $m_1 = \{CID_i, G_{ij}, H_{ij}, Z_i, Ni\}$

### Verification Procedure:

$R_i = G_{ij} \text{ XOR } h(h(y) || Ni || SID_j) = 11 \text{ XOR } h(28 || 6 || 7) = 11 \text{ XOR } 29 = 22$

$M_i = h(R_i || h(x || y)) = h(22 || 23) = 24$

$L_i = H_{ij} \text{ XOR } h(M_i || Ni || SID_j) = 63 \text{ XOR } h(24 || 6 || 7) = 63 \text{ XOR } 38 = 25$

$h(r \text{ XOR } PWi) = CID_i \text{ XOR } h(L_i || M_i || Ni) \rightarrow 21=48 \text{ XOR } h(25 || 24 || 6) = 48 \text{ XOR } 37 = 21$

$E_i = M_i \text{ XOR } h(r \text{ XOR } PWi) = 24 \text{ XOR } 21 = 13$

$h(E_i || M_i || Ni) ?= Z_i \rightarrow h(13 || 24 || 6) = 32 \rightarrow 32 = 32 \rightarrow$  Yes

Generates  $N_j = 8$

$V_{ij} = h(E_i || Ni || M_i || SID_j) = h(13 || 6 || 24 || 7) = 33$

Second parameter Set  $m_2 = \{V_{ij}, N_j\}$

$h(E_i || Ni || M_i || SID_j) ?= V_{ij} \rightarrow h(13 || 6 || 24 || 7) = 33 \rightarrow 33 = 33 \rightarrow$  Yes

$V_{ij}' = h(E_i || N_j || M_i || SID_j) = h(13 || 8 || 24 || 7) = 34$

Third Parameter Set  $m_3 = \{V_{ij}'\}$

$h(E_i || N_j || M_i || SID_j) ?= V_{ij}' \rightarrow h(13 || 8 || 24 || 7) = 34 \rightarrow 34=34 \rightarrow$  Yes

$SK = h(E_i || Ni || N_j || M_i || SID_j) = h(13 || 6 || 8 || 24 || 7) = 35$

### Password Change Procedure:

$L_i = W_i \text{ XOR } h(ID_i || h(r \text{ XOR } PWi)) = 61 \text{ XOR } h(1 || 21) = 61 \text{ XOR } 36 = 25$

$F_i^* = h(L_i) = h(25) = 27$

$F_i^* ?= F_i \rightarrow 27 = 27 \rightarrow$  Yes

Select  $PW_{new}=9$ ,  $r_{new}=10$

$W_{new} = L_i \text{ XOR } h(ID_i || h(r_{new} \text{ XOR } PW_{new})) = 25 \text{ XOR } h(1 || h(10 \text{ XOR } 9)) = 25 \text{ XOR } h(1 || h(3)) = 25 \text{ XOR } h(1 || 36) = 25 \text{ XOR } 39 = 62$

$E_{new} = E_i \text{ XOR } h(r \text{ XOR } PWi) \text{ XOR } h(r_{new} \text{ XOR } PW_{new}) = 13 \text{ XOR } 21 \text{ XOR } 36 = 60.$

All the phases are mentioned above with the numerical analysis, it shows how the user enters his credential, how it will verify, the overall scenario of scheme cleared with the numerical analysis.

### 3. REVIEW OF CHEN’S SCHME

**Table 1 - Computation Cost Analysis**

Phases Schemes	Registration phase	Login phase	Authentication phase	Password Change phase	Total
Chen’s Scheme	$10Th + 6Tx + 4TC$	$11Th + 9Tx + 9TC$	$14Th + 5Tx + 30TC$	$7Th + 6Tx + 2TC$	$42Th + 26Tx + 45Tc$

**Table 2- Entity Wise Analysis**

Entities Schemes	User	Server	Registration Center	Total
Chen’s Scheme	$23Th + 16Tx + 21TC$	$12Th + 5Tx + 20TC$	$8Th + 5Tx + 4TC$	$42Th + 26Tx + 45Tc$

Where,

Th = Time for one way hash function

Tx = Time for XOR function

Tc = Time for Concat function

### 4. CONCLUSION

This paper analyzes the Chen’s Scheme in mathematical terms. Chen’s Scheme is based on two factor Remote user authentication scheme. We analyzed the functions performed in each and every phase of authentication scheme and also analyzed functions performed by every participating entity. In future, we will try to make our proposed scheme which is more robust compared to Chen’s scheme.

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